EE 360C — Algorithms — Summer 2013
Homework #3

Due: July 10, 2013 11:30am (in class)

Homework problems are to be done individually. You may discuss the problem and general concepts with other students, but you must write your solutions independently.

Each question is worth 10 points. Maximum possible score is 30.

Whenever you give an algorithm, prove that it is correct.

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1. Suppose you're a consultant for the networking company CluNet, and they have the following problem. The network that they're currently working on is modeled by a connected graph $G = (V, E)$ with n nodes. Each edge e is a fiber-optic cable that is owned by one of two companies—creatively named X and Y—and leased to CluNet.

   Their plan is to choose a spanning tree T of G and upgrade the links corresponding to the edges of T. Their business relations people have already concluded an agreement with companies X and Y stipulating a number k so that in the tree T that is chosen, k of the edges will be owned by X and $n - k - 1$ of the edges will be owned by Y.

   CluNet management now faces the following problem. It is not at all clear to them whether there even *exists* a spanning tree T meeting these conditions, or how to find one if it exists. So this is the problem they put to you: Give a polynomial-time algorithm that takes G, with each edge labeled X or Y, and either (i) returns a spanning tree with exactly k edges labeled X, or (ii) reports correctly that no such tree exists

2. One of the basic motivations behind the Minimum Spanning Tree Problem is the goal of designing a spanning network for a set of nodes with minimum total cost. Here we explore another type of objective: designing a spanning network for which the most expensive edge is as cheap as possible.

   Specifically, let $G = (V, E)$ be a connected graph with n vertices, m edges, and positive edge costs that you may assume are all distinct. Let $T = (V, E')$ be a spanning tree of G; we define the bottleneck edge of T to be the edge of T with the greatest cost.

   A spanning tree T of G is a minimum-bottleneck spanning tree if there is no spanning tree T' of G with a cheaper bottleneck edge.

   (a) Is every minimum-bottleneck tree of G a minimum spanning tree of G? Prove or give a counterexample.

   (b) Is every minimum spanning tree of G a minimum-bottleneck tree of G? Prove or give a counterexample.

3. Let $G = (V, E)$ be an (undirected) graph with costs $c_e \geq 0$ on the edges $e \in E$. Assume you are given a minimum-cost spanning tree T in G. Now assume that a new edge is added to G, connecting two nodes $v, w \in V$ with cost c.

   (a) Give an efficient algorithm to test if T remains the minimum-cost spanning tree with the new edge added to G (but not to the tree T). Make your algorithm run in time $O(|E|)$. Can you do it in $O(|V|)$ time? Please note any assumptions you make about what data structure is used to represent the tree T and the graph G.

   (b) Suppose T is no longer the minimum-cost spanning tree. Give a linear-time algorithm (time $O(|E|)$) to update the tree T to the new minimum-cost spanning tree.